# Advanced Parallel Programming

Miscellaneous MPI-IO topics

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# **MPI-IO Errors**

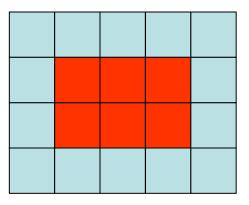


- Unlike the rest of MPI, MPI-IO errors are not fatal
  - probably don't want your program to crash if a file open fails
  - always need to check the error code!
- Many different error codes can be reported
  - I would suggest simply quitting if ierr != MPI\_SUCCESS
- Can change this behaviour for file operations
  - same functionality as MPI\_Errhandler\_create etc.
  - called MPI\_File\_create\_errhandler, ...
  - error handlers are attached to file handles rather than communicators
  - can set handler to be MPI\_ERRORS\_ARE\_FATAL

- Useful to check length of output file
  - ls -l <filename>
  - check that size (in bytes) is what you expect
- Can be confusing if file already exists
  - length will be increased if new file is longer than existing file
  - but may not be decreased if new file is shorter!
- Delete old files before running your test programs

#### Datatype for MPI File read/write

- Usually pass the basic type of the array being processed
  - eg MPI\_FLOAT, MPI\_REAL
- Can pass derived types
  - useful for receiving the core of an array when local arrays have halos

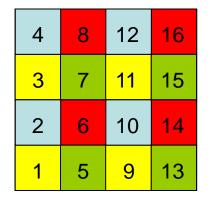


MPI\_File\_read\_all(fh, &x[1][1], 1, vector3x2, ...);
MPI\_FILE\_READ\_ALL(fh, x(2,2) , 1, vector3x2, ...)

– or could use a 3x2 subarray and pass &x[0][0] or x(1,1)

# **General Decompositions**

- We have just considered block decompositions
  - where local array size is an exact multiple of global array size
- If the sizes don't match
  - define different sized subarrays on each process
  - eg processes at the edge of the grid have smaller subsections
- This does not generalize to block-cyclic decompositions
  - how do we specify discontinuous subarrays?





#### **Distributed Arrays**

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int MPI\_Type\_create\_darray(int size, int rank, int ndims, int array\_of\_gsizes[], int array\_of\_distribs[], int array\_of\_dargs[], int array\_of\_psizes[], int order, MPI Datatype oldtype, MPI Datatype \*newtype);

MPI\_TYPE\_CREATE\_DARRAY(SIZE, RANK, NDIMS
ARRAY\_OF\_GSIZES, ARRAY\_OF\_DISTRIBS, ARRAY\_OF\_DARGS,
ARRAY\_OF\_PSIZES, ORDER, OLDTYPE, NEWTYPE, IERR)

INTEGER SIZE, RANK, NDIMS, ARRAY\_OF\_GSIZES(\*), ARRAY\_OF\_DISTRIBS(\*), ARRAY\_OF\_DARGS(\*), ARRAY\_OF\_PSIZES(\*), ORDER, OLDTYPE, NEWTYPE, IERR

- See the man page for full details!
  - uses HPF conventions for block-cyclic distributions

#### **Unstructured Data**

- Imagine a particle simulation
  - each particle is a compound object with a type and position (x,y,z)
    - eg a C struct or Fortran type
  - each particle has unique global identifier 1, 2, 3, ..., N-1, N
- Particles move around
  - at the end of a simulation, each process will have:
    - a different number of particles
    - with a random mixture of global identifiers
- Two choices
  - write to file in the order they appear in the processes
  - write to file with position based on global identifier

### Approach

- |epcc|
- Define a derived type to match the particle object
  - eg MPI\_PARTICLE
  - use this as the etype
- Writing in process order
  - need to know where to start in the file
  - calculate the sum of the number of particles on previous ranks
    - using MPI\_Scan
- Writing in global order
  - call MPI\_Type\_indexed (Or create\_indexed\_block)
  - use this as the filetype
  - write multiple instances of MPI\_PARTICLE

#### **Unstructured Meshes**

- Similar to global ordering of particles
  - each element has both a local and global identifier
  - want the file to be ordered by the global id
- Define an MPI\_ELEMENT
  - use this as the etype
  - create an indexed filetype based on global id



This code spends a lot of time waiting while saving to disk

define big arrays: old and new

loop many times

- ! do a computationally expensive operation
- new = expensive\_function(old)

```
old = new
```

every 10 iterations:

```
save to disk(old)
```

end loop

# Non-blocking IO

• This code overlaps computation and IO

define big arrays: old and new

```
loop many times
```

```
! do a computationally expensive operation
new = expensive function(old)
if (saving to disk):
  finish: isave to disk(old)
old = new
every 10 iterations:
  start: isave to disk(old)
```

end loop

# Non-blocking IO in MPI-IO

- Two forms
- General non-blocking
  - MPI\_File\_iwrite(fh, buf, count, datatype, request)
  - finish by waiting on **request**
  - but no collective version
- Split collective
  - MPI\_File\_write\_all\_begin(fh, buf, count, datatype)
  - MPI\_File\_write\_all\_end(fh, buf, status)
  - only a single outstanding IO operation at any one time
  - allows for collective version

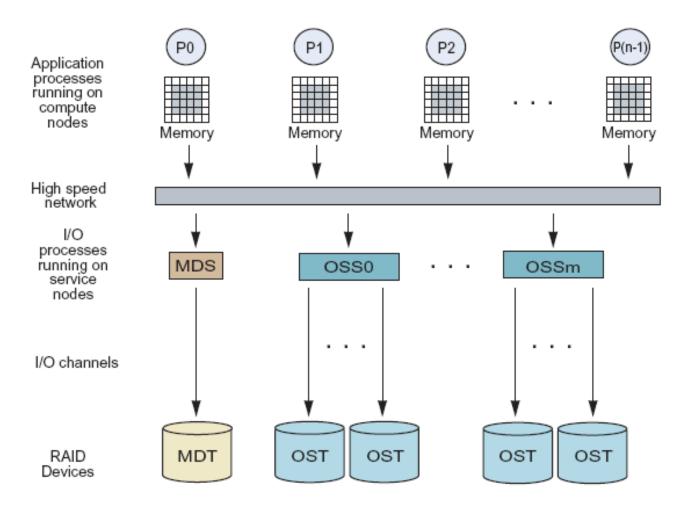


- How can I read MPI-IO files in a serial program?
- Using native format
  - data is raw bytes
  - use fread in C or direct access unformatted IO in Fortran
  - see ioread.c and ioread.f90 for examples
  - Fortran approach is quite old-fashioned (direct access IO)
     new access="stream" functionality makes this a bit simpler
- Other MPI-IO formats will require more work!
- Note that you can do single process IO in MPI-IO
  - pass MPI\_COMM\_SELF to MPI\_File\_open

#### • I have advised

- define a datatype to represents mapping from local to global data
- use this in MPI\_File\_set\_view()
- then do linear reads / writes; gaps are automatically skipped
- Alternative approach
  - let everyone see the whole file (i.e. do not set a view)
  - manually seek to correct location using, e.g., MPI\_File\_write\_at()
  - displacement is in units of the extent of the etype
- Disadvantages
  - a very low-level, manual approach less amenable to IO optimisation
  - danger that each request is handled individually with no aggregation
  - can use MPI\_File\_write\_at\_all() but might still be slow

• Recall schematic overview of parallel file system Lustre

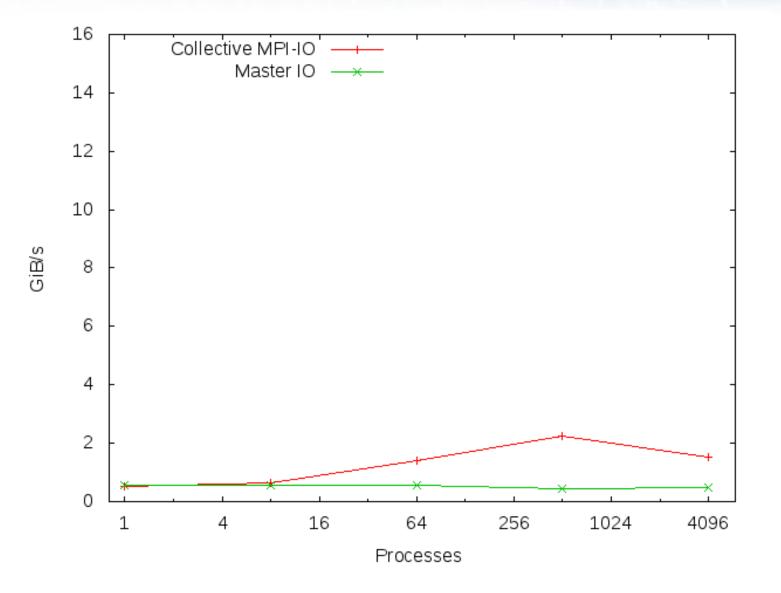


MPI-IO 4: Miscellaneous

# **Application-side parallel IO**

- Implementing MPI-IO has achieved
  - all data going to a single file
  - minimal stress on Meta Data Server (MDS) a serial bottleneck
  - potential for many processes to write simultaneously
- But ...
  - performance requires multiple parallel writes to disk
  - in Lustre, requires multiple Object Storage Servers (OSS) writing to multiple Object Storage Targets (OST)
  - an OSS is like an IO server, an OST is like a physical disk
- User has control over assignment of files to OSTs
  - but default is only a few OSTs
  - MPI-IO performance not much better than naïve master IO

### Parallel vs serial IO, default Lustre



#### MPI-IO 4: Miscellaneous

#### **Cellular Automaton Model**

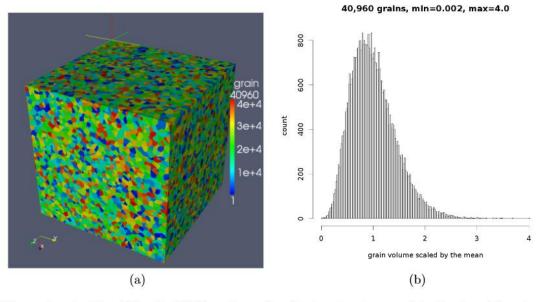


Figure 1: A  $4.1 \times 10^9$  cell, 40,960 grain equiaxed microstructure model, showing (a) grain arrangement with colour denoting orientation; (b) grain size size (volume) histogram.

 Fortran coarray library for 3D cellular automata microstructure simulation, Anton Shterenlikht, proceedings of 7<sup>th</sup> International Conference on PGAS Programming Models, 3-4 October 2013, Edinburgh, UK.

#### Benchmark

- Distributed regular 3D dataset across 3D process grid
  - local data has halos of depth 1; set up for weak scaling
  - implemented in Fortran and MPI-IO

! Write data collectively

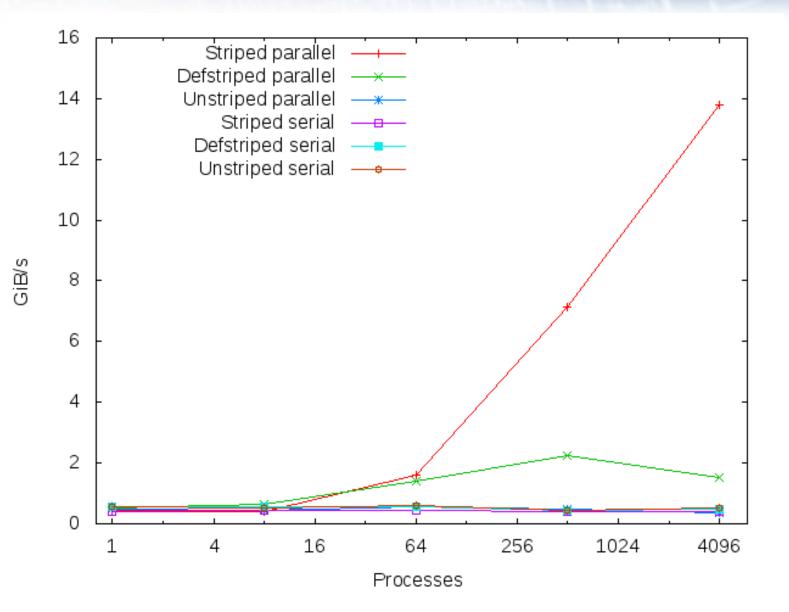
call MPI\_File\_write\_all(fh, iodata, 1, mpi\_subarray, status, ierr)

- Can split a file across multiple OSTs
  - each block is called a "stripe"; default striping is across 4 OSTs

# • lfs setstripe -c 8 <directory>

- stripes across 8 OSTs for all files in the directory
- has substantial benefits for performance
- stripe count of "-1" means use all OSTs
- Test case
  - 128 x 128 x 128 array of doubles on each process in 3D grid
  - scaled up to 4096 processes = 64 GiB
  - identical IO approach as used in exercise
    - generalised to 3D
    - local halos automatically stripped off with derived type in MPI-IO write call

#### **Results on ARCHER**





# **Performance Summary**

- Serial IO never gets more than about 500 MiB/s
  - peak for a single OST
- With default striping, never exceed 2 GiB/s
  - 4 stripes = 4 OSTs = 4 x 500 MiB/s
- With full striping, IO bandwith increases with process count
  - can achieve in excess of 10 GiB/s
- Collective IO is essential
  - replacing MPI\_File\_Write\_all()
     by MPI\_File\_write() disastrous!
  - identical functionality but each IO request now processed separately with file locking

Processes	Bandwidth
1	49.5 MiB/s
8	5.9 MiB/s
64	2.4 MiB/s



- MPI web pages
- Short ARCHER report:
  - http://www.archer.ac.uk/documentation/white-papers/
- Another tutorial
  - https://www.lrde.epita.fr/~ricou/mpi-io.ppt
- Advanced MPI book
  - "Using Advanced MPI: Modern Features of the Message-Passing Interface"

