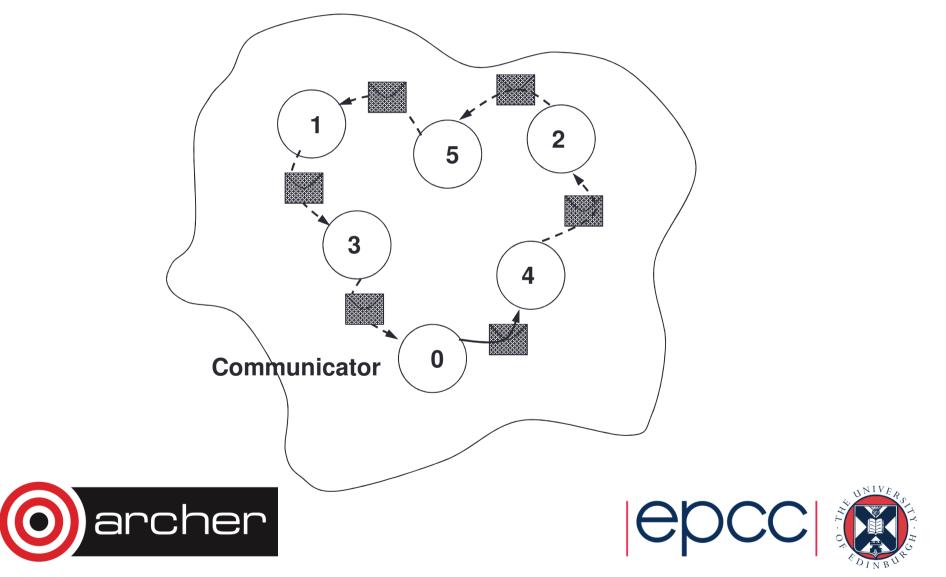
Non-Blocking Communications



Deadlock



Completion

- The *mode* of a communication determines when its constituent operations complete.
 - i.e. synchronous / asynchronous
- The *form* of an operation determines when the procedure implementing that operation will return
 - i.e. when control is returned to the user program





Blocking Operations

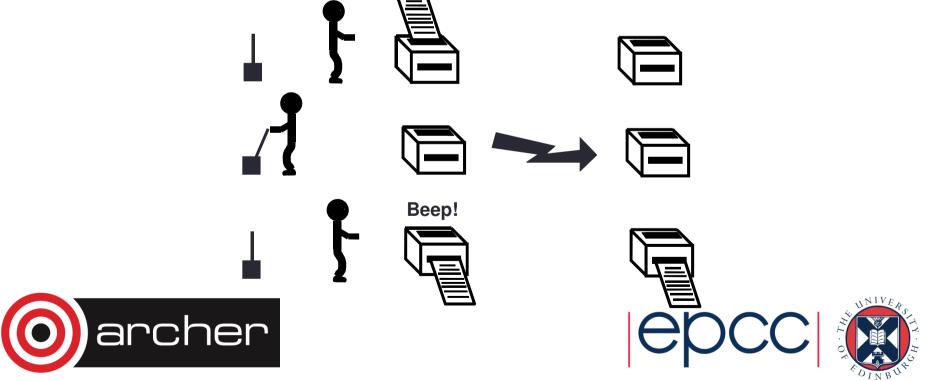
- Relate to when the operation has completed.
- Only return from the subroutine call when the operation has completed.
- These are the routines you used thus far
 - MPI_Ssend
 - MPI_Recv





Non-Blocking Operations

 Return straight away and allow the sub-program to continue to perform other work. At some later time the sub-program can *test* or *wait* for the completion of the non-blocking operation.



Non-Blocking Operations

- All non-blocking operations should have matching wait operations. Some systems cannot free resources until wait has been called.
- A non-blocking operation immediately followed by a matching wait is equivalent to a blocking operation.
- Non-blocking operations are not the same as sequential subroutine calls as the operation continues after the call has returned.





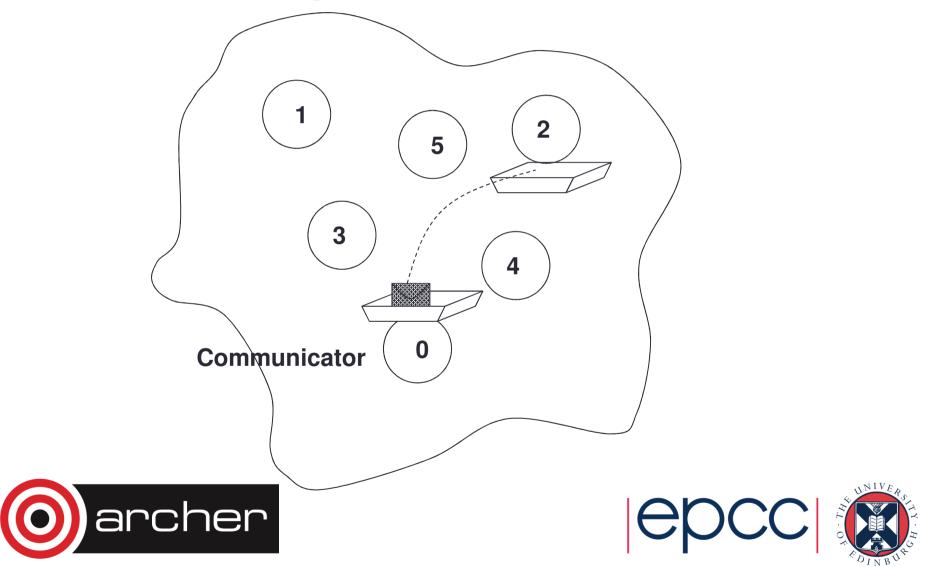
Non-Blocking Communications

- Separate communication into three phases:
- Initiate non-blocking communication.
- Do some work (perhaps involving other communications?)
- Wait for non-blocking communication to complete.

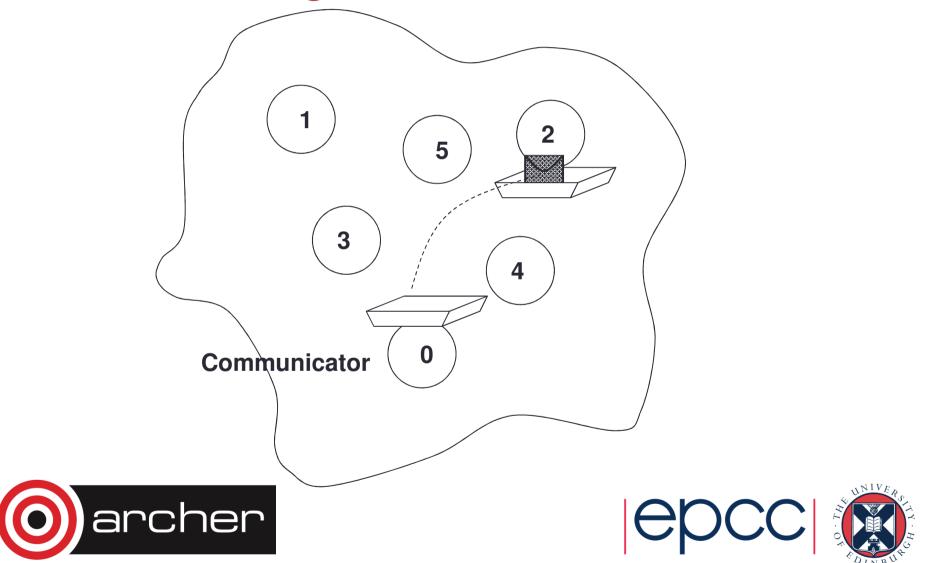




Non-Blocking Send



Non-Blocking Receive



Handles used for Non-blocking Comms

- datatype same as for blocking (MPI_Datatype or INTEGER).
- communicator same as for blocking (MPI_Comm or INTEGER).
- request MPI_Request or INTEGER.
- A request handle is allocated when a communication is initiated.





Non-blocking Synchronous Send

• C:

• Fortran:

MPI_WAIT(request, status, ierror)





Non-blocking Receive

• C:

• Fortran:

MPI_WAIT(request, status, ierror)





Send data from rank 1 to rank 3

```
! Array of ten integers
integer, dimension(10) :: x
integer :: reqnum
integer, dimension(MPI_STATUS_SIZE) :: status
.....
```

```
if (rank .eq. 1)
CALL MPI_WAIT(reqnum, status, ierr)
```





Blocking and Non-Blocking

- Send and receive can be blocking or non-blocking.
- A blocking send can be used with a non-blocking receive, and vice-versa.
- Non-blocking sends can use any mode synchronous, buffered, standard, or ready.
- Synchronous mode affects completion, not initiation.





Communication Modes

NON-BLOCKING OPERATION	MPI CALL
Standard send	MPI_ISEND
Synchronous send	MPI_ISSEND
Buffered send	MPI_IBSEND
Ready send	MPI_IRSEND
Receive	MPI_IRECV





Completion

• Waiting versus Testing.

• C:

• Fortran:

MPI_WAIT(handle, status, ierror)

MPI_TEST(handle, flag, status, ierror)





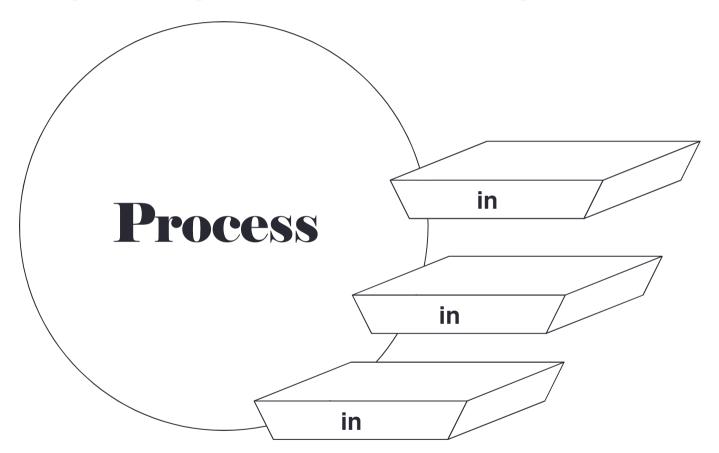
Multiple Communications

- Test or wait for completion of one message.
- Test or wait for completion of all messages.
- Test or wait for completion of as many messages as possible.





Testing Multiple Non-Blocking Comms







Combined Send and Receive

Specify all send / receive arguments in one call

- MPI implementation avoids deadlock
- useful in simple pairwise communications patterns, but not as generally applicable as non-blocking





Exercise

Rotating information around a ring

- See Exercise 4 on the sheet
- Arrange processes to communicate round a ring.
- Each process stores a copy of its rank in an integer variable.
- Each process communicates this value to its right neighbour, and receives a value from its left neighbour.
- Each process computes the sum of all the values received.
- Repeat for the number of processes involved and print out the sum stored at each process.





Possible solutions

- Non-blocking send to forward neighbour
 - blocking receive from backward neighbour
 - wait for forward send to complete
- Non-blocking receive from backward neighbour
 - blocking send to forward neighbour
 - wait for backward receive to complete
- Non-blocking send to forward neighbour
- Non-blocking receive from backward neighbour
 - wait for forward send to complete
 - wait for backward receive to complete





Notes

- Your neighbours do not change
 - send to left, receive from right, send to left, receive from right, ...
- You do not alter the data you receive
 - receive it
 - add it to you running total
 - pass the data unchanged along the ring
- You must not access send or receive buffers until communications are complete
 - cannot read from a receive buffer until after a wait on irecv
 - cannot overwrite a send buffer until after a wait on issend



