## **Cray I/O Software Enhancements**

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#### **Overview**

- The Cray Linux Environment and parallel libraries provide full support for common I/O standards.
  - Serial POSIX I/O
  - Parallel MPI I/O
    - 3<sup>rd</sup> part-libraries built on top of MPI I/O
    - HDF5, NetCDF4
- Cray versions provide many enhancements over generic implementations that integrate directly with Cray XC30 and Cray Sonexion hardware.
  - Cray MPI-IO collective buffering, aggregation and data sieving.
  - Automatic buffering and direct I/O for Posix transfers via IOBUF.
- This talk explains how to get the best from the enhanced capabilities of the Cray software stack.

# **Cray MPI-IO Layer**

**Data Aggregation and Sieving** 

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- The MPI-2.0 standard provides a standardised interface for reading and writing data to disk in parallel. Commonly referred to as MPI I/O
- Full integration with other parts of the MPI standard allows users to use derived type to complete complex tasks with relative ease.
- Can automatically handle portability like byte-ordering and native and standardised data formats.
- Available as part of the cray-mpich library on XC30, commonly referred to as Cray MPI-IO.
  - Fully optimised and integrated with underlying Lustre file-system.

#### **Step 1: MPI-IO Hints**

The MPI I/O interface provides a mechanism for providing additional information about how to the MPI-IO layer should access files.

These are controlled via MPI-IO HINTS, either via calls in the MPI API or passed via an environment variable. All hints can be set on a file-by-file basis.

On the Cray XC30 the first most useful are:

• striping\_factor – Number of lustre stripes

 striping\_unit – Size of lustre stripes in bytes
 These set the file's Lustre properties when it is created by an MPI-IO API call.

\* Note these require MPICH\_MPIIO\_CB\_ALIGN to be set to its default value of 2.

#### **Example settings Lustre hints in C**

Hints can be added to MPI calls via an Info unit when the file is opened using the MPI I/O API. Below is an example in C

```
#include <mpi.h>
#include <stdio.h>
int factor = 4; // The number of stripes
int unit = 4; // The stripe size in megabytes
sprintf(factor_string, "%d", factor);
// Multiple unit into bytes from megabytes
sprintf(unit string, "%d", unit * 1024 * 1024);
MPI_Info_set(info, "striping_factor", factor_string);
MPI Info set(info, "striping unit", unit string);
MPI File open(MPI COMM WORLD, filename, MPI MODE CREATE
              MPI MODE RDWR, info, &fh);
```

## **Setting hints via environment variables**

Alternatively, hints can be passed externally via an environment variable, MPICH\_MPIIO\_HINTS.

Hints can be applied to all files, specific files, or pattern files, e.g.

# Set all MPI-IO files to 4 x 4m stripes
MPICH\_MPIIO\_HINTS="\*:striping\_factor=4:striping\_unit=4194304"

# Set all .dat files to 8 x 1m stripes
MPICH\_MPIIO\_HINTS="\*.dat:striping\_factor=8:striping\_unit=1048576"

#### **Displaying hints**

The MPI-IO library can print out the "hint" values that are being using by each file when it is opened. This is controlled by setting the environment variable:

export MPICH\_MPIIO\_HINT\_DISPLAY=1

The reported is generated by the PE with rank 0 in the relevant communicator and is printed to stderr.

```
PE 0: MPICH/MPIIO environment settings:
       MPICH MPIIO HINTS DISPLAY = 1
PE 0:
       MPICH MPIIO HINTS
PE 0:
                                 = NULL
PE 0:
       MPICH MPIIO_ABORT_ON_RW_ERROR = disable
PE 0:
       MPICH MPIIO CB ALIGN
                               = 2
PE 0:
       MPIIO hints for file1:
•••
         direct io
                                 = false
         aggregator placement stride = -1
•••
```

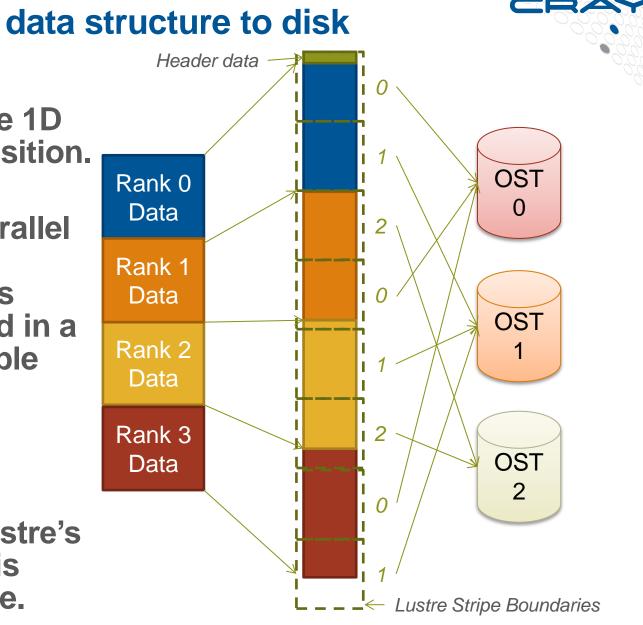
#### **Collective vs independent calls**

- Opening a file via MPI I/O is a collective operation that must be performed by all members of a supplied communicator.
- However, many individual file operations have two versions:
  - A collective version which must be performed by all members of the supplied communicator
  - An independent version which can be performed ad-hoc by any processor at any time. This is akin to standard POSIX I/O, however includes MPI data handling syntactic sugar.
- It is only during collective calls that the MPI-IO library can perform required optimisations. Independent I/O is usually no more (or less) efficient than POSIX equivalents.

# **Collective Buffering & Data Sieving**

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#### Writing a simple data structure to disk

**Consider a simple 1D** parallel decomposition.

**MPI I/O allows parallel** data structures distributed across ranks to be stored in a single with a simple offset mapping.

However exactly matching this distribution to Lustre's stripe alignment is difficult to achieve.

#### **Recap: Optimising Lustre Performance**

Lustre's performance comes from Parallelism, with many writers/readers to/from many Object Storage Targets (OSTs).

MPI I/O offers good parallelism, with each rank able potentially writing it's own data into a file

However, for large jobs #writers >> #OSTs, and each rank may write to more than 1 OST. This can cause Lustre lock contention and that slows access

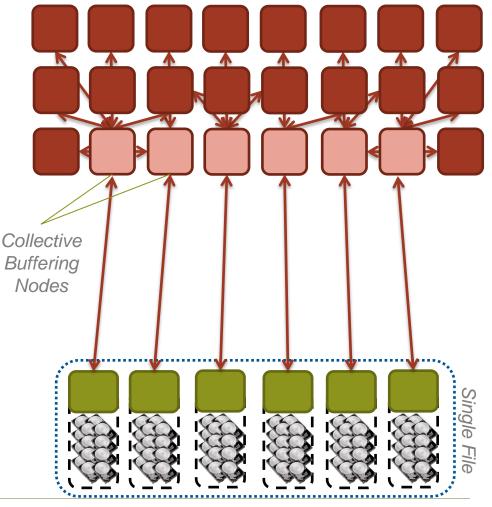
**W** 14 4 K Lock Points

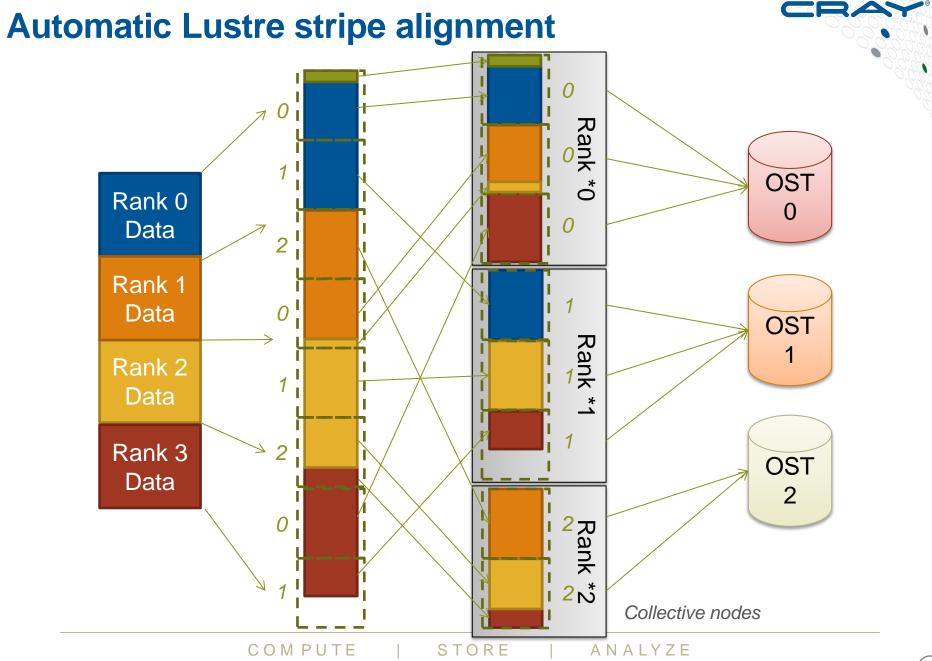
#### **Collective Buffering and Lustre Stripe Alignment**

To limit the number of writers the MPI-IO library will assign and automatically redistribute data to a subset of "collective buffering" or "aggregator" nodes during a collective file operation.

By default, the number of "collective buffering" nodes will be the same as the lustre striping factor to get maximum benefit of Lustre Stripe Alignment.

Each collective buffer node will attempt to only write data to a single Lustre OST.

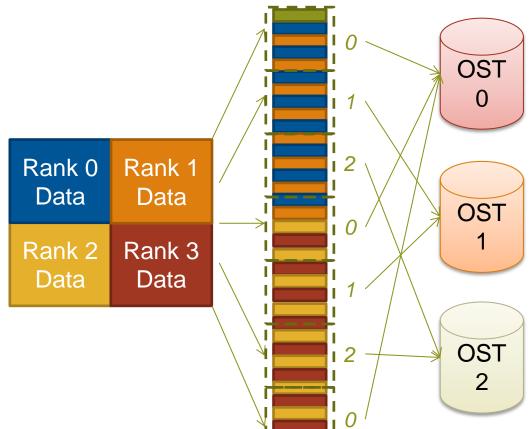




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#### Writing structured data to disk

However, switching to an even slightly more complex decomposition, like a 2D Cartesian, results in ranks having to perform noncontiguous file operations.





#### **Data Sieving**

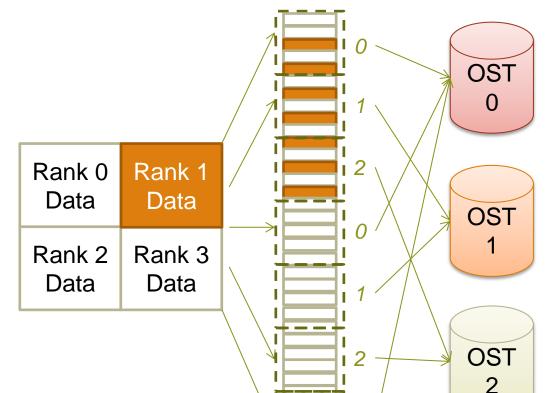
- "Read/Write Gaps" occur when the data is not accessed contiguously from the file.
- This limits the total bandwidth rate as each access requires separate calls and may cause additional seek time on HDD storage.
- Overall performance can be improved by minimising the number of read/write gaps.
- The Cray MPI-IO library will attempt to use data sieving to automatically combine multiple smaller operations into fewer larger operations.

#### **Strided file access**

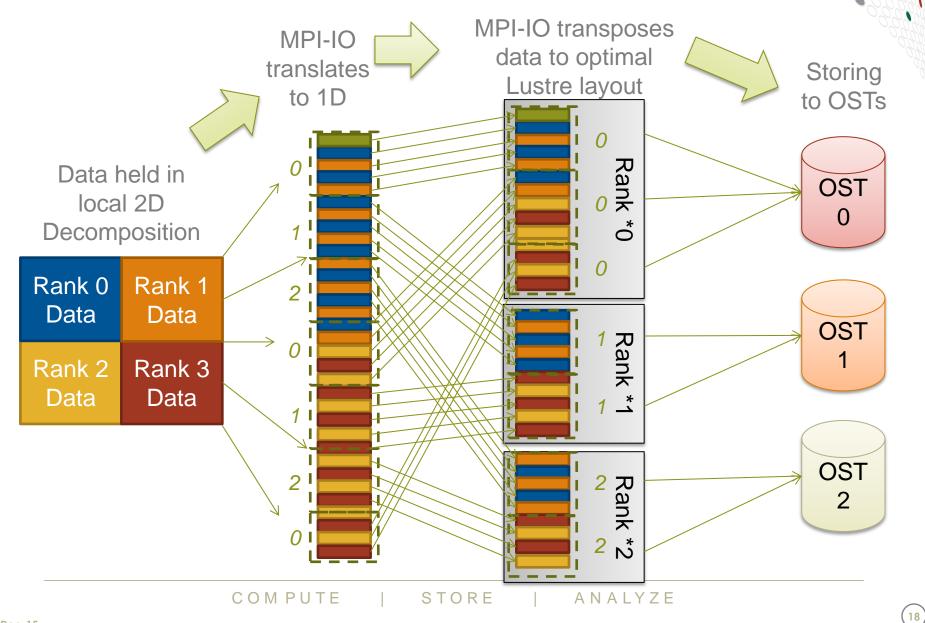
Focusing on a rank we can see that it will potentially end up writing strided data to each OST.

This is likely to incur penalties due to extent locking on each of the OSTs.

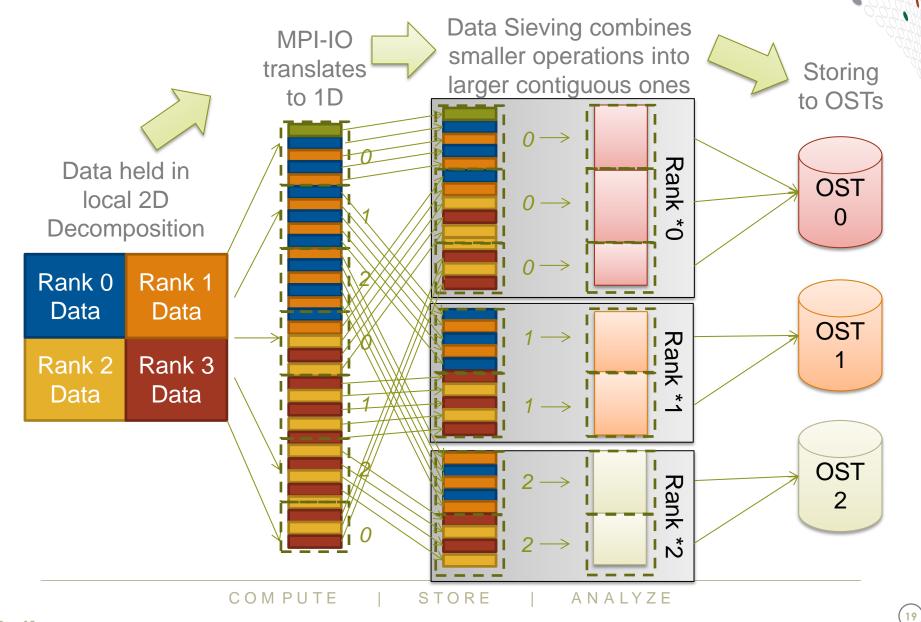
It also prevents optimal performance of HDD block devices writing contiguous blocks of data



#### Writing structured data to disk



#### **Data Sieving**



## **Managing Collective Buffering**

- The Cray MPI-IO library will automatically perform collective buffering of collective MPI-IO calls. There are two algorithms controlled by the value of MPICH\_MPIIO\_CB\_ALIGN=[0|2]
  - 0 : distribute data into equally across all aggregators regardless of Lustre strip settings (inefficient if data in a single stripe or small number of stripes)
  - 2 (default): Divides data into Lustre stripe-sized pieces and assigns them to collective buffering nodes such that each node always and exclusively accesses the same set of stripes.

#### • The default behaviour (MPICH\_MPIIO\_CB\_ALIGN=2) will:

- Automatically set the number of aggregators to the number of stripes
- Attempt to place each aggregator on it's own node

• So in most cases it is only necessary to change the Lustre stripe settings to optimise performance

### Manually configuring collective buffer

- It is possible to specify specific values for collective buffering. This is done by setting MPICH\_MPIIO\_CB\_ALIGN=0
- The primary tunable values of collective buffering are:
  - cb\_buffers\_size The size of each buffer (default 16MB)
  - cb\_nodes The number of aggregators (default # of XC30 nodes)
- The are passed as hints to the MPI-IO library
- Other variables are explained in man mpi
- Our experiences is that the default aligned algorithm achieves best performance in most circumstances.

#### **Understanding MPI-IO Stats**

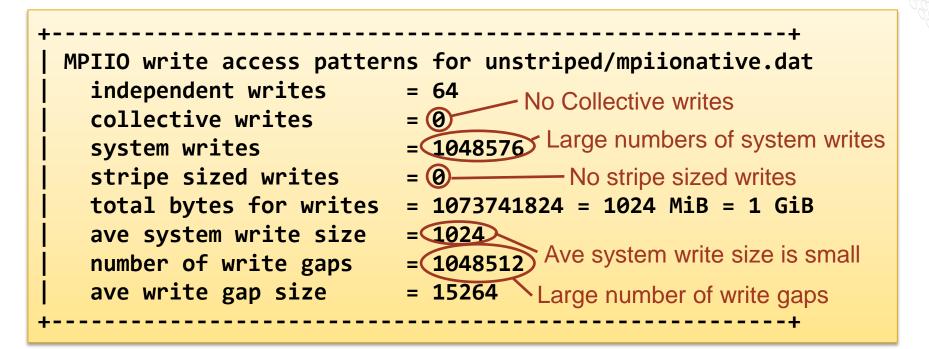
The MPI library can provide stats on how many reads and writes were performed in system sized gaps. Adding: export MPICH\_MPII0\_STATS=1 To run time environment variables will generate summary output on each PE.

+	+				
MPIIO write access patterns for file1					
independent writes	= 0				
collective writes	= 24				
system writes	= 4916				
<pre>stripe sized writes</pre>	= 4915				
<pre>total bytes for writes</pre>	= 25769803776 = 24576 MiB = 24 GiB				
ave system write size	= 5242026				
number of write gaps	= 0				
ave write gap size	= NA				
+	+				

#### In more detail

- Independent writes the number of writes performed by independent call to the MPI-IO library
- Collective writes the number of writes performed in collective MPI-IO calls.
- System writes the number of POSIX write operations the MPI-IO translated the calls into
- Total bytes for writes The amount of data written to the file
- Avg system write size The average size of each POSIX write operation
- Number of write gaps the number of gaps/seeks between POSIX write operations
- Avg write gap size the average size of jumps/seek operations.

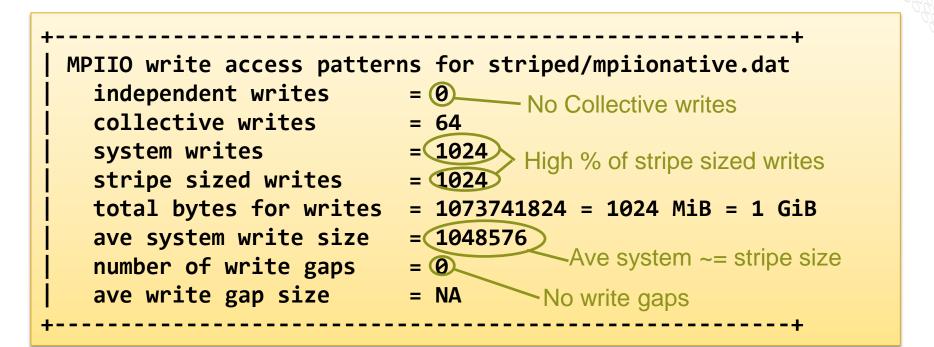
#### **Recognising Poor Performance**



This is a simple example for 3D decomposed array. Independent MPI-IO writes are used in place of collectives.

## 0.005 GiB/s

#### **Recognising Good Performance**



This same simple example for 3D decomposed array. Now using collective MPI-IO writes:

## 1.41 GiB/s

#### **Cray-mpich 7 features**



#### export MPICH\_MPIIO\_AGGREGATOR\_PLACEMENT\_DISPLAY=1

#### Aggregator Placement for /lus/scratch/myfile

RankReorderMethod=3 AggPlacementStride=-1

AGG	Rank	nid	
0	0	nid00578	
1	4	nid00579	
2	1	nid00606	
3	5	nid00607	
4	2	nid00578	
5	6	nid00579	
6	3	nid00606	
7	7	nid00607	

# **Controlling IO Buffering**

In traditional serial IO

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## **IO Buffering**

 By default the underlying Linux OS tries to automatically buffer all IO

- The user does not have any control over the buffering process.
- The OS tries to use as much of the free memory as possible as buffers
  - There are limits that can be set in place by the admins of the system
  - These are not controlled by the user
- Users can request to skip OS buffering by using 'direct IO', however this requires modifying the open system call (0\_DIRECT)

#### Cray offers a more sophisticated IO buffer method named IOBUF

- Available via a module in the Cray PE and controlled via a runtime Environment Variable
- User can control buffer sizes for each file
- Will automatically pre-fetch data
- Provides summary statistics
- Man page available

#### IOBUF

- IOBUF is a library that intercepts standard I/O (stdio) and enables asynchronous caching and prefetching of sequential file access
- Should not be used for
  - Hybrid programs that do I/O within a parallel region (not thread-safe)
  - Many processes accessing the same file in a coordinated fashion (MPI\_File\_write/read\_all)

#### • No need to modify the source code but just

- Load the module iobuf
- Relink your application
- Set export IOBUF\_PARAMS='\*:verbose' in the batch script
- See the iobuf(3) manpage

### **IOBUF STATS output**

IOBUF paramete file="defstrip :verbose	rs: ed/serial.d	at":size=104857	6:count=4:vbuf	fer_count=4096:pi	refetch=1		
PE 0: File "defstriped/serial.dat"							
	Calls	Seconds	Megabytes	Megabytes/sec	Avg Size		
Write	2048	0.580566	402.653184	693.552615	196608		
Open	1	0.001288					
Close	1	0.006056					
Buffer Write	384	0.533518	402.653184	754.713968	1048576		
I/O Wait	384	0.530056	402.653184	759.643408			
Buffers used	4	(4 MB)					
Preflushes	384						

- Each file accessed on each PE will print a summary when closed.
- Users set a "buffer size" (default 1MB), transactions that are smaller are cached into one of the buffers
- Larger transactions are performed directly, bypassing the buffers.

#### **IOBUF** configuration

- Users can increase the size of buffers (size=#[KMG])
- They can also add more buffers (count=#) this allows for more access points
- Data is automatically pre-fetched. More buffers can be prefetched (count=#) or disabled completely (count=0)
- Data can also be written "direct", i.e. bypassing the OS's internal buffering process.
- Settings controlled on a file by file basis or via pattern matching, e.g:

#### Conclusions

- The Cray XC30 has a fully featured and optimised MPI I/O stack available to users – If all possible – Use it!
  - It is fully integrated with the underlying Lustre file system
  - Also integrated with supplied HDF5 and NetCDF4 libraries.
- Application I/O should be parallel with multiple writers to achieve best performance.
- Lustre settings should match application parallelism
  - (e.g. file-per-process vs MPI-IO collectives)
  - Lustre settings have biggest impact on perfromance.
- MPI-IO provides nice way to abstract complicated fileaccess patterns
  - But implementations can only optimise parallel collective operations in practice!