



# Parallel Models

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Different ways to exploit parallelism

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# Outline

- Shared-Variables Parallelism
  - threads
  - shared-memory architectures
- Message-Passing Parallelism
  - processes
  - distributed-memory architectures
- Practicalities
  - compilers
  - libraries
  - usage on real HPC architectures



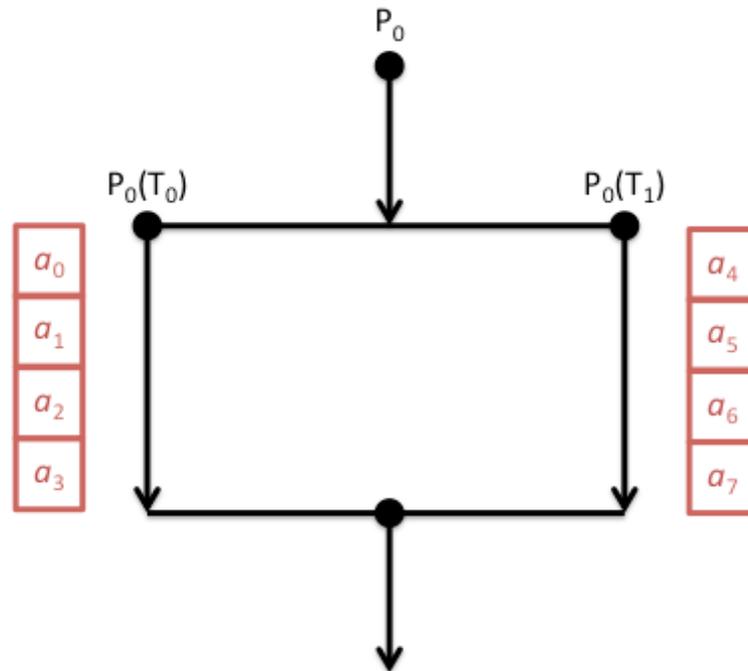
# Shared Variables

Threads-based parallelism



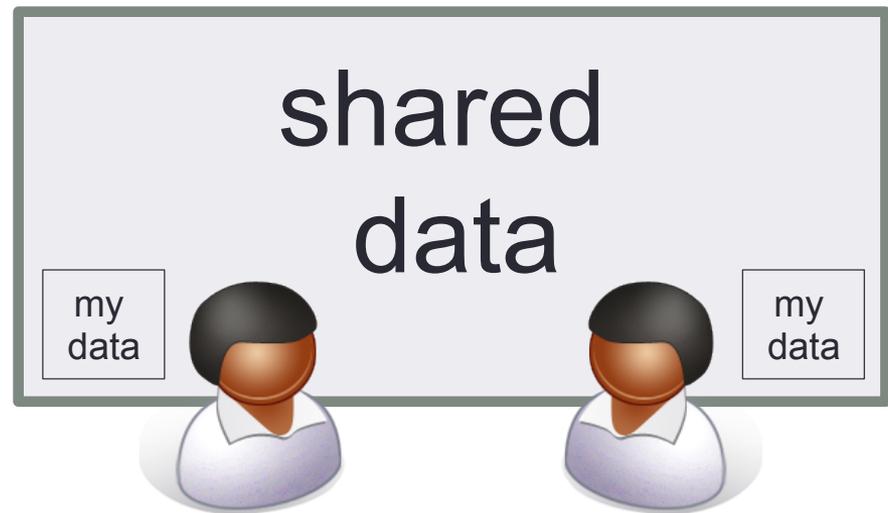
# Shared-memory concepts

- Have already covered basic concepts
  - threads can all see data of parent process
  - can run on different cores
  - potential for parallel speedup



# Analogy

- One very large whiteboard in a two-person office
  - the shared memory
- Two people working on the same problem
  - the threads running on different cores attached to the memory
- How do they collaborate?
  - working together
  - but not interfering
- Also need *private* data



# Thread Communication

Thread 1

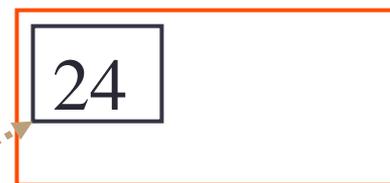
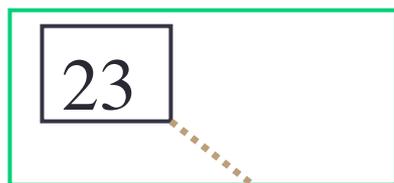
Thread 2

Program

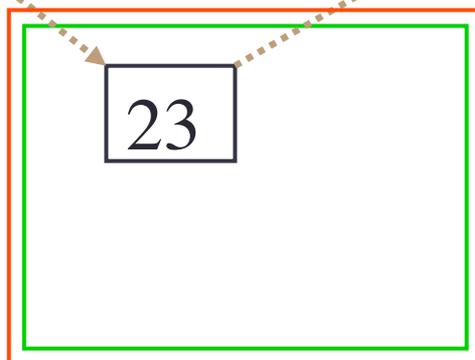
```
mya=23  
a=mya
```

```
mya=a+1
```

Private  
data



Shared  
data



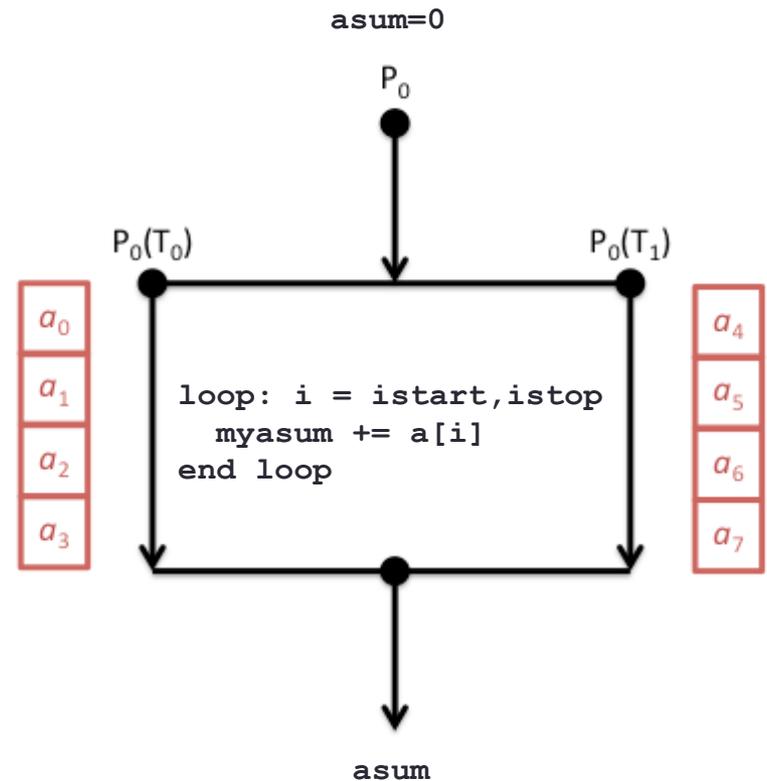
# Synchronisation

- Synchronisation crucial for shared variables approach
  - thread 2's code must execute *after* thread 1
- Most commonly use global barrier synchronisation
  - other mechanisms such as locks also available
- Writing parallel codes relatively straightforward
  - access shared data as and when its needed
- Getting correct code can be difficult!



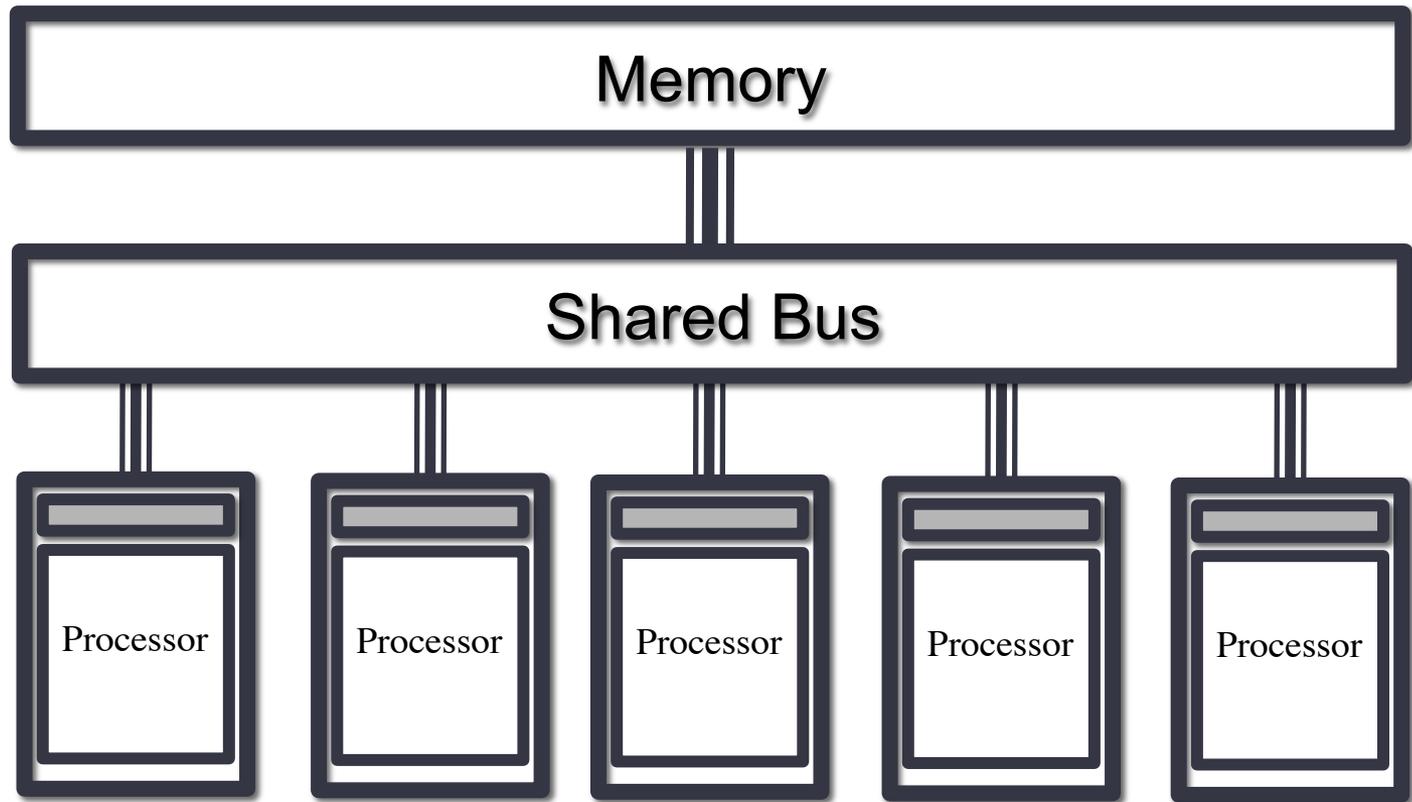
# Specific example

- Computing  $asum = a_0 + a_1 + \dots + a_7$ 
  - shared:
    - main array: `a[8]`
    - result: `asum`
  - private:
    - loop counter: `i`
    - loop limits: `istart, istop`
    - local sum: `myasum`
  - synchronisation:
    - thread0: `asum += myasum`
    - barrier
    - thread1: `asum += myasum`



# Hardware

- Needs support of a shared-memory architecture



# Threads: Summary

- Shared blackboard a good analogy for thread parallelism
- Requires a shared-memory architecture
  - in HPC terms, cannot scale beyond a single node
- Threads operate independently on the shared data
  - need to ensure they don't interfere; synchronisation is crucial
- Threading in HPC usually uses OpenMP directives
  - supports common parallel patterns
  - e.g. loop limits computed by the compiler
  - e.g. summing values across threads done automatically



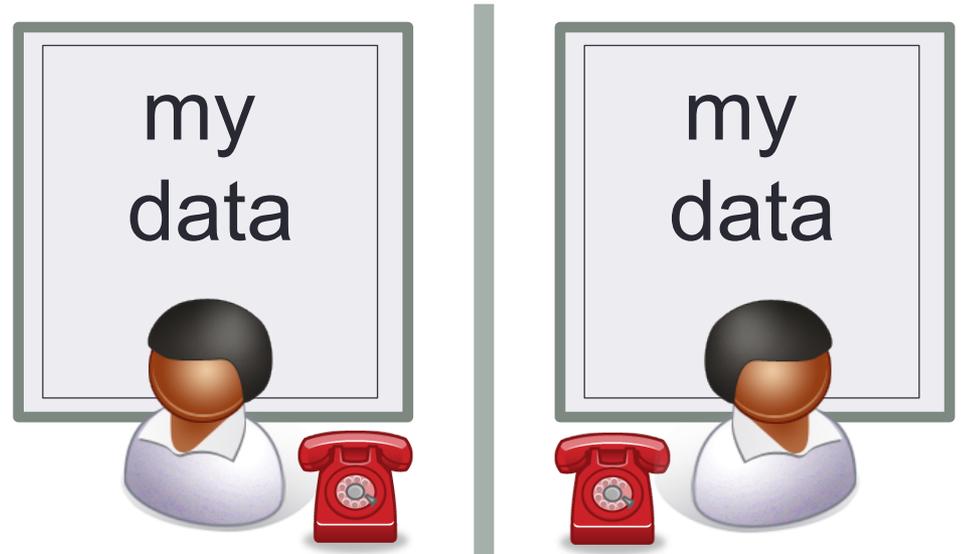
# Message Passing

Process-based parallelism



# Analogy

- Two whiteboards in different single-person offices
  - the distributed memory
- Two people working on the same problem
  - the processes on different nodes attached to the interconnect
- How do they collaborate?
  - to work on single problem
- Explicit communication
  - e.g. by telephone
  - no shared data



# Process communication

## Process 1

Program

$a=23$

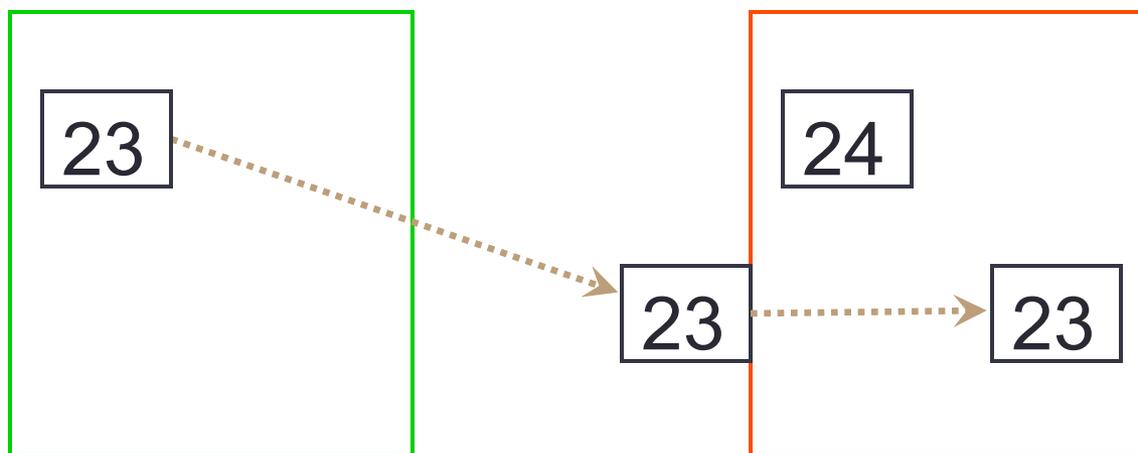
$\text{Send}(2, a)$

## Process 2

$\text{Recv}(1, b)$

$a=b+1$

Data

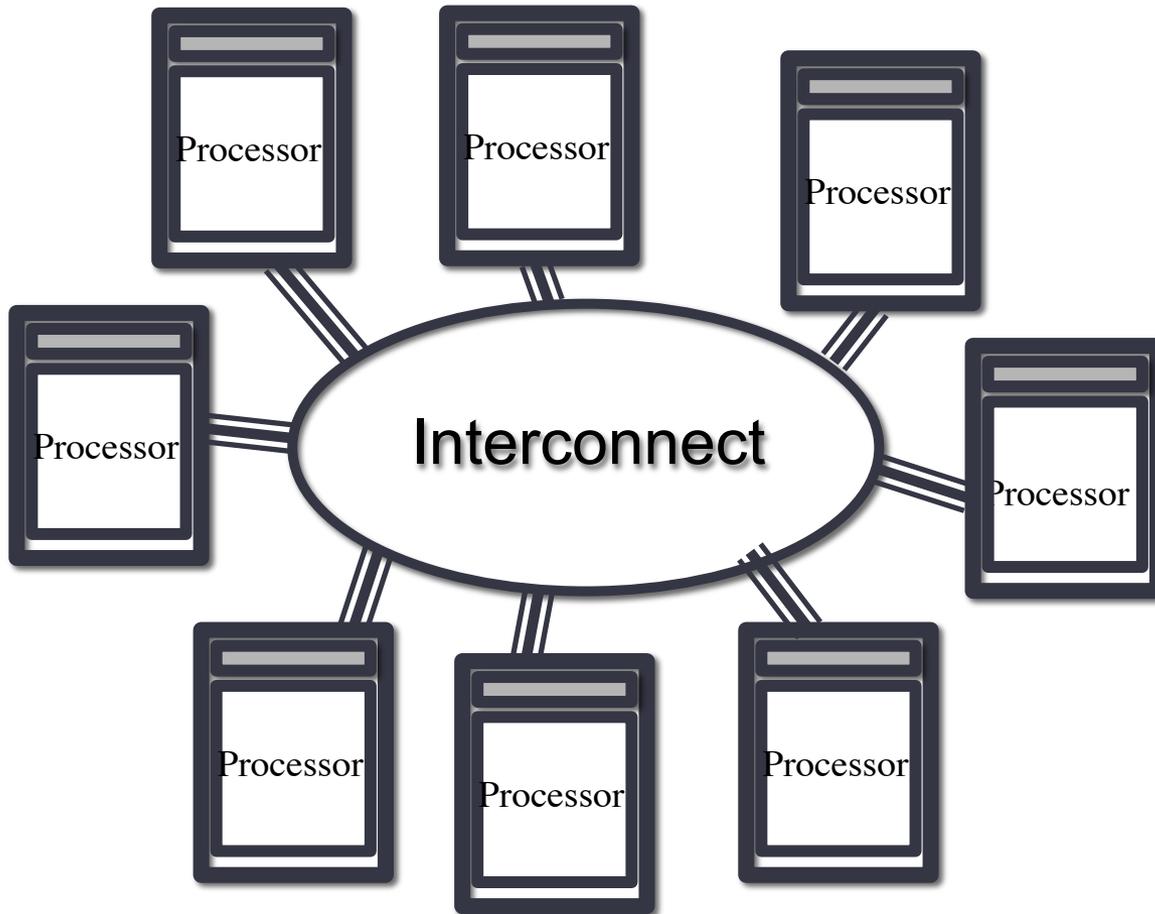


# Synchronisation

- Synchronisation is automatic in message-passing
  - the messages do it for you
- Make a phone call ...
  - ... wait until the receiver picks up
- Receive a phone call
  - ... wait until the phone rings
- No danger of corrupting someone else's data
  - no shared blackboard



# Hardware



- Natural map to distributed-memory
  - one process per processor-core
  - messages go over the interconnect, between nodes/OS's

# Processes: Summary

- Processes cannot share memory
  - ring-fenced from each other
  - analogous to white boards in separate offices
- Communication requires explicit *messages*
  - analogous to making a phone call, sending an email, ...
  - synchronisation is done by the messages
- Almost exclusively use Message-Passing Interface
  - MPI is a library of function calls / subroutines



# Practicalities



Interconnect

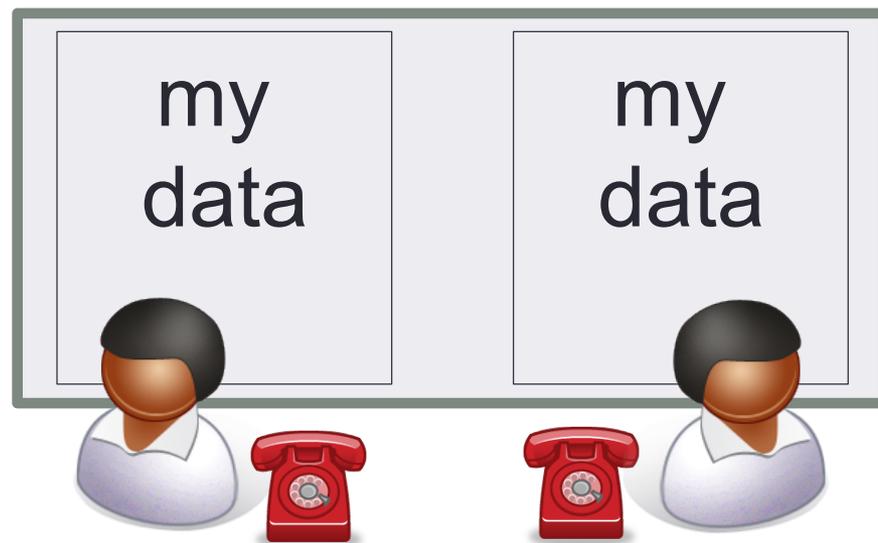


- 8-core machine might only have 2 nodes
  - how do we run MPI on a real HPC machine?
- Mostly ignore architecture
  - pretend we have single-core nodes
  - one MPI process per processor-core
  - e.g. run 8 processes on the 2 nodes
- Messages between processes on the same node are fast
  - but remember they also share access to the network

# Message Passing on Shared Memory

- Run one process per core
  - don't directly exploit shared memory
  - analogy is phoning your office mate
  - actually works well in practice!

- Message-passing programs run by a special job launcher
  - user specifies #copies
  - some control over allocation to nodes



# Summary

- Shared-variables parallelism
  - uses threads
  - requires shared-memory machine
  - easy to implement but limited scalability
  - in HPC, done using OpenMP compilers
- Distributed memory
  - uses processes
  - can run on any machine: messages can go over the interconnect
  - harder to implement but better scalability
  - on HPC, done using the MPI library

