Advanced Parallel Programming

Overview of Parallel IO

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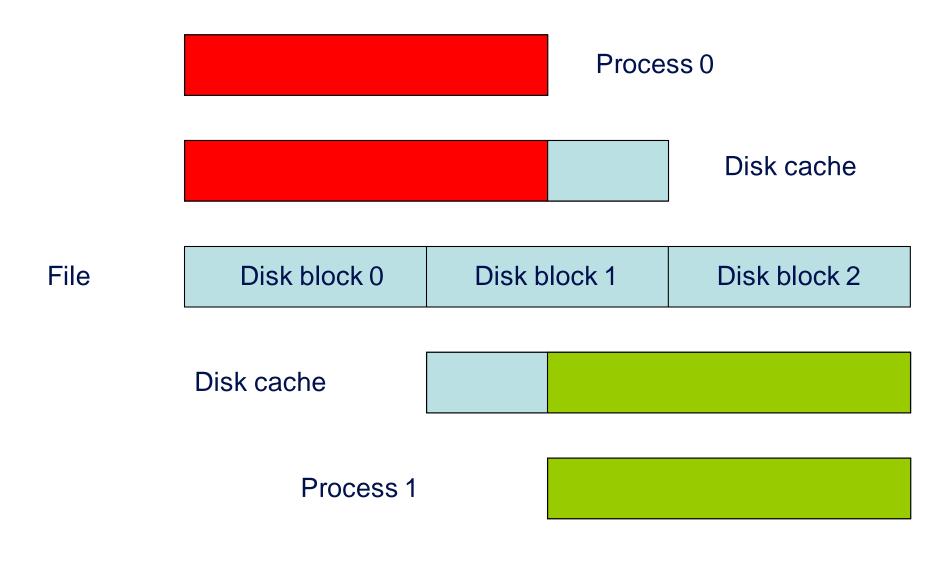
- Lecture will cover
 - Why is IO difficult
 - Why is parallel IO even worse
 - Straightforward solutions in parallel
 - What is parallel IO trying to achieve?
 - Files as arrays
 - MPI-IO and derived data types



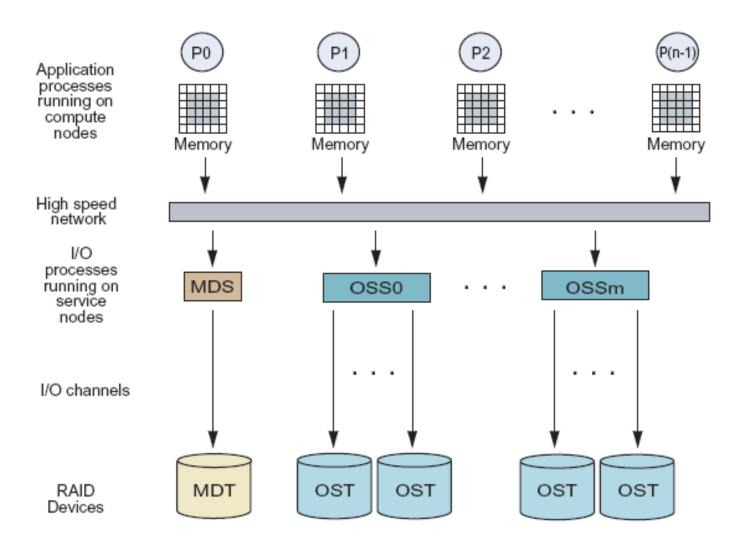
- Breaks out of the nice process/memory model
 - data in memory has to physically appear on an external device
- Files are very restrictive
 - linear access probably implies remapping of program data
 - just a string of bytes with no memory of their meaning
- Many, many system-specific options to IO calls
- Different formats
 - text, binary, big/little endian, Fortran unformatted, ...
- Disk systems are very complicated
 - RAID disks, many layers of caching on disk, in memory, ...
- IO is the HPC equivalent of printing!

- Cannot have multiple processes writing a single file
 - Unix generally cannot cope with this
 - data cached in units of disk blocks (eg 4K) and is not coherent
 - not even sufficient to have processes writing to distinct parts of file
- Even reading can be difficult
 - 1024 processes opening a file can overload the filesystem (fs)
- Data is distributed across different processes
 - processes do not in general own contiguous chunks of the file
 - cannot easily do linear writes
 - local data may have halos to be stripped off

Simultaneous Access to Files



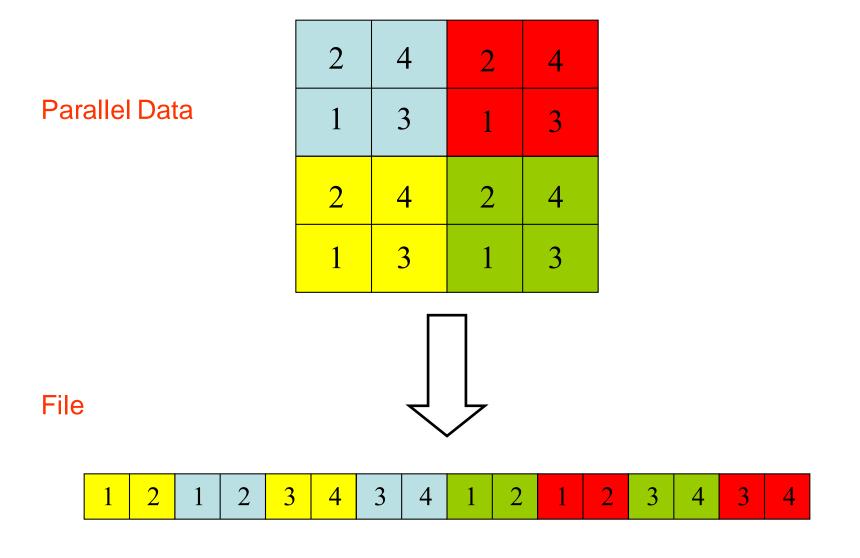
Parallel File Systems: Lustre



Parallel File Systems

- Allow multiple IO processes to access same file
 - increases bandwidth
- Typically optimised for bandwidth
 - not for latency
 - e.g. reading/writing small amounts of data is very inefficient
- Very difficult for general user to configure and use
 - need some kind of higher level abstraction
 - focus on data layout across user processes
 - don't want to worry about how file is split across IO servers

4x4 array on 2x2 Process Grid





- imagine a shared array called \mathbf{x}

begin serial region
 open the file
 write x to the file
 close the file
end serial region

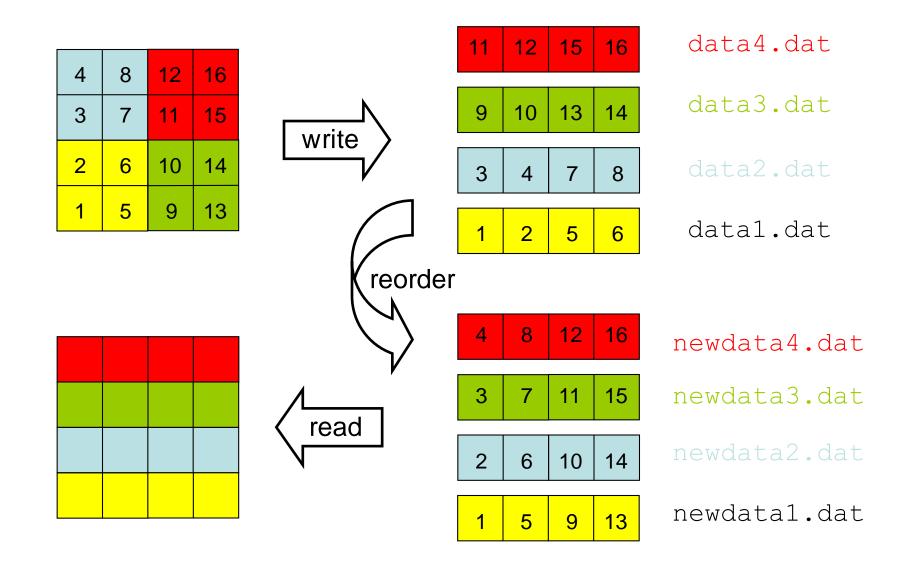
Simple as every thread can access shared data
 may not be efficient but it works

• But what about message-passing?

Message Passing: Naive Solutions

- Master IO
 - send all data to/from master and write/read a single file
 - quickly run out of memory on the master
 - or have to write in many small chunks
 - does not benefit from a parallel fs that supports multiple write streams
- Separate files
 - each process writes to a local fs and user copies back to home
 - or each process opens a unique file (dataXXX.dat) on shared fs
- Major problem with separate files is reassembling data
 - file contents dependent on number of CPUs and decomposition
 - pre / post-processing steps needed to change number of processes
 - but at least this approach means that reads and writes are in parallel
 - but may overload filesystem for many processes

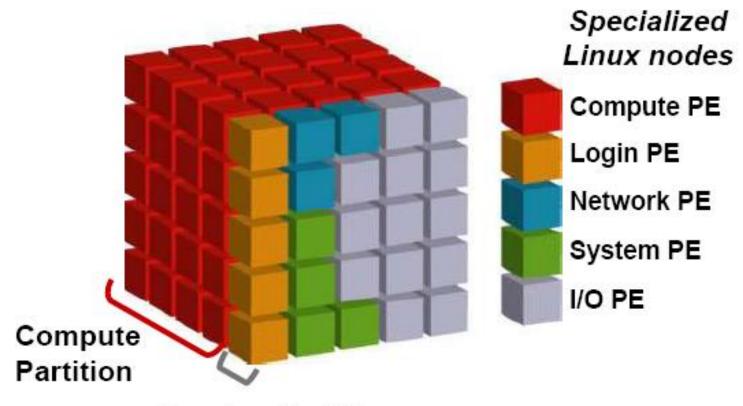
2x2 to 1x4 Redistribution



- A way to do parallel IO properly
 - where the IO system deals with all the system specifics
- Want a single file format
 - We already have one: the serial format
- All files should have same format as a serial file
 - entries stored according to position in global array
 - not dependent on which process owns them
 - order should always be 1, 2, 3, 4,, 15, 16

- What does the IO system need to know about the parallel machine?
 - all the system-specific fs details
 - block sizes, number of IO servers, etc.
- All this detail should be hidden from the user
 - but the user may still wish to pass system-specific options ...

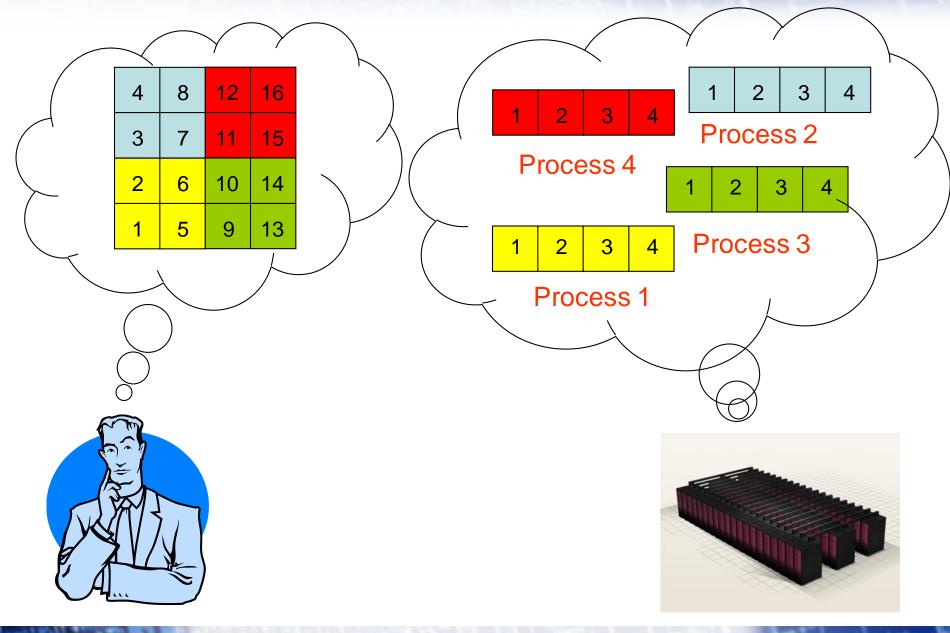
Example of IO system: Cray XT4



Service Partition

- What does the IO system need to know about the data?
 - how the local arrays should be stitched together to form the file
- But ...
 - mapping from local data to the global file is only in the mind of the programmer!
 - the program does not know that we imagine the processes to be arranged in a 2D grid
- How do we describe data layout to the IO system
 - without introducing a whole new concept to MPI?
 - cartesian topologies are not sufficient
 - do not distinguish between block and block-cyclic decompositions

Programmer View vs Machine View



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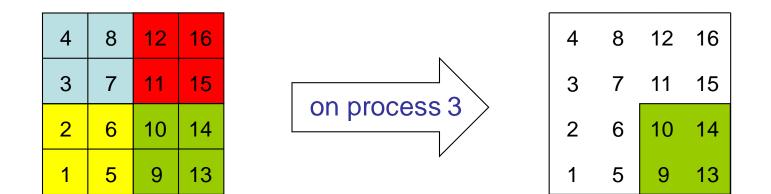
MPI-IO 1: Overview of Parallel IO

Files vs Arrays

- Think of the file as a large array
 - forget that IO actually goes to disk
 - imagine we are recreating a single large array on a master process
- The IO system must create this array and save to disk
 - without running out of memory
 - never actually creating the entire array
 - ie without doing naive master IO
 - and by doing a small number of large IO operations
 - merge data to write large contiguous sections at a time
 - utilising any parallel features
 - doing multiple simultaneous writes if there are multiple IO nodes
 - managing any coherency issues re file blocks

MPI-IO Approach

- MPI-IO is part of the MPI-2 standard
 - http://www.mpi-forum.org/docs/docs.html
- Each process needs to describe what subsection of the global array it holds
 - it is entirely up to the programmer to ensure that these do not overlap for write operations!
- Programmer needs to be able to pass system-specific information
 - pass an info object to all calls



- Describe 2x2 subsection of 4x4 array
- Using standard MPI derived datatypes
- A number of different ways to do this
 - we will cover three methods in the course

- Parallel IO is difficult
 - in theory and in practice
- MPI-IO provides a high-level abstraction
 - user describes global data layout using derived datatypes
 - MPI-IO hides all the system specific fs details …
 - ... but (hopefully) takes advantage of them for performance
- User requires a good understanding of derived datatypes
 see next lecture