

Parallel IO in Code_Saturne

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Code_Saturne Main Features and Toolchain Two Applications

Motivation Code_Saturne IO Methods On the Fly Mesh Generation: Mesh Multiplication

Test Architectures and Test Cases Scalability at Scale I/O using HECToR (Lustre) Results - ARCHER (Lustre) vs Blue Joule (GPFS) Conclusions - Perspectives





- Code_Saturne is developed by EDF (France)
- Computational Fluid Dynamics
- open source
- Fortran, C, Python
- fully validated production versions with long-term support every two years (currently 3.0)
- development versions
- http://code-saturne.org



Code_Saturne's Features

Technology

-Co-located finite volume, arbitrary unstructured meshes, predictor-corrector -350 000 lines of code, 37% Fortran, 50% C, 13% Python

-MPI for distributed-memory and some openMP for shared-memory machines

Physical modeling

- -Laminar and turbulent flows: k-eps, k-omega, SST, v2f, RSM, LES models -Radiative transfer (DOM, P-1)
- -Coal, heavy-fuel and gas combustion
- -Electric arcs and Joule effect
- -Lagrangian module for particles tracking
- -Atmospheric modeling (merging Mercure_Saturne)
- -ALE method for deformable meshes
- -Rotor / stator interaction for pump modeling, for marine turbines

Flexibility

- -Portability (Unix, Linux and MacOS X)
- -Graphical User Interface with possible integration within the SALOME platform



Toolchain

Reduced number of tools

- Each with rich functionality
- Natural separation between interactive and potentially long-running parts
- In-line (pdf) documentation





Example Applications

Free surface modelling (ALE)



Thermofluids study of the hot box dome AGR (EDF Energy)



- Complex flow due through the forest of tubes
- Calculation shows little mixing in the centre of the dome
- Temperatures at the dome highest where thermocouples are located



Code_Saturne I/O

Different types of file I/O

- read input
- write checkpoint data periodically
- read checkpoint if restarting a previous simulation
- write output

Different methods for I/O

- STD C IO
- MPI IO



Motivation (1)

High-End Machines offer hope for more multi-physics & multi-scale for engineering in ever more detailed configurations.

Huge effort has been dedicated to improve/optimise solvers (in our case Navier-Stokes equation solvers) for them to scale on the current existing petaflop machines, but arguably less time is dedicated by CFD developers to IOs.

Several types of IOs and some way around loading/writing huge data files have been identified:-

-INPUT: mesh, domain partition (if already known), restart file (if needed), input data

-OUTPUT: mesh (if changed, with added periodicity for instance), domain partition (if computed by the code), listing file, postprocessing file, checkpoint, probes



Ways around exist to avoid loading full data set for:--INPUT:-

-mesh (mesh joining and mesh multiplication)

-domain partition (partition re-computed by the code) -OUTPUT:-

-pre-processed mesh (not needed, because computed by the code)

-domain partition (not needed because computed by the code)

-post-processing (co-processing, for instance using Catalyst)

But not for:-

-INPUT:-

-restart file, as/if the whole flow field is needed -OUTPUT:-

-checkpoint file, as/if the whole flow field is needed



I/O Methods in 3.3.1

I/O Method

CS_FILE_STDIO_SERIAL

Serial standard C IO (funnelled through rank 0 in parallel)

CS_FILE_STDIO_PARALLEL

CS_FILE_MPI_INDEPENDENT

Per-process standard C IO

Non-collective MPI-IO with independent file open and close

CS_FILE_MPI_NON_COLLECTIVE

Non-collective MPI-IO with collective file open and close

CS_FILE_MPI_COLLECTIVE

Collective MPI-IO



I/O Methods in 3.3.1

Selecting the I/O method

• GUI and XML file

- -> "Calculation Management"
 - -> "Performance Tuning"

• Directly:

- Can be set in the cs_user_performance_tuning file in cs_user_parallel_io()
- Can also provide MPI IO hints





Use global numbering Redistribution on n blocks

- n blocks ≤ n cores
- Minimum block size may be set to avoid many small blocks (for some communication or usage schemes), or to force 1 block (for I/O with non-parallel libraries)

21	22	23	24	25
16	17	18	19	20
11	12	13	14	15
6	7	8	9	10
1	2	3	4	5

Rank 0 is collecting info from the blocks





Mesh Multiplication

Most mesh generators are serial and thus memory-limited

A way around to generate extremely large meshes is to build meshes from existing coarse ones and globally refine each cell

This process might be repeated several times

Developed by Ales Ronovsky (VSB, PRACE)



(a)







(c)



Architectures

ARCHER – XC30 / Lustre

3008 Compute nodes: two 2.7 GHz, 12-core E5-2697 v2 (Ivy Bridge) series processors. Within the node, QuickPath Interconnect (QPI) links to connect the 2 processors

The Cray Aries interconnect links all compute nodes in a Dragonfly topology.

Compute nodes access the file system via IO nodes running the Cray Data Virtualization Service (DVS)

Blue Joule – BGQ / GPFS

6 racks, each rack containing 1,024 16-core, 64 bit, 1.60 GHz A2 PowerPC processors.

All the racks have 8 IO nodes which connect the BGQ racks to the shared GPFS storage over Infiniband.

The minimum block size which can be booted for a job is therefore 1,024/8 nodes, or 128 nodes.



Test Case - Configuration

3D lid-driven cavity - fully unstructured mesh (tetras)

Size of the meshes:

- MM Level 0 (13 million cells Current production runs)
- MM Level 1 (111 million cells Current production runs)
- MM Level 2 (890 million cells Production runs in 2015)
- MM Level 3 (7.2 billion cells Production runs in 2016/2017)

Geometric partitioning using a Space-Filling Curve approach (Hilbert)



IO tests are performed when the solver performance is still acceptable

If not stated, machine default settings. No striping for Lustre, for instance



Scalability at Scale (1)

Mesh generated by Mesh Multiplication



105B Cell Mesh (MIRA, BGQ)

Cores	Time in Solver
262,144	652.59s
524,288	354.89s

13B Cell Mesh (MIRA, BGQ)

Nodes/Ranks	Time in Solver	
16384/32	70.124s	
32768/32	50.207s	
49152/32	43.465s	





Scalability at Scale (2)

Comparison HECToR – ARCHER

Mesh generated by Mesh Multiplication Cube meshed with tetra cells





IO HECToR (Lustre)

Comparison IO per Blocks (Ser-IO) and MPI-IO Comparison Lustre (Cray) / GPFS (IBM BlueGene/Q)

	HECToR		Blue Joule	
Cores	MPI-IO	Ser-IO	MPI-IO	Ser-IO
2048	633	1203	-	-
4096	608	640	85	1279
8192	859	1147	86	1300
16384	732	747	67	1330
32768	-	-	59	1360

Tube Bundle

812M cells



Block IO: ~same performance on Lustre and GPFS MPI-IO: 8 to 10 times faster with GPFS



















Quick Summary





MPI-IO - 7.2 B Tetra Mesh



MPI – IO vs Block IO

Writing Checkpoint Files – Mesh_Output







Conclusions

With the current machine/filesystem settings

MPI-IO

ARCHER (Lustre) better for small meshes than larger ones BlueJoule (GPFS) better for large meshes than smaller ones

MPI-IO vs Block IO

If results on HECToR were comparable, much better obtained with MPI-IO on ARCHER





Lustre and Striping

Previous ARCHER results used defaults for striping. Use striping for better performance for large meshes?

Stripe count for results directory set to all available OSTs with:

Ifs setstripe



Striping – MM Level1

MPI-IO - 111 M Tetra Mesh



No Stripping Read Input 814MB No Stripping Write Checkpoint1 1.7GB No Stripping Write Checkpoint2 3.3GB No Stripping Write Mesh_Output 11.6GB Full Stripping Read Input 814MB Full Stripping Write Checkpoint1 1.7GB Full Stripping Write Checkpoint2 3.3GB Full Stripping Write Mesh_Output 11.6GB



Striping – MM Level 2

MPI-IO - 890 M Tetra Mesh



No Stripping Read Input 814MB No Stripping Write Checkpoint1 13.5GB No Stripping Write Checkpoint2 26.5GB No Stripping Write Mesh_Output 92.8GB Full Stripping Read Input 814MB Full Stripping Write Checkpoint1 13.5GB Full Stripping Write Checkpoint2 26.5GB Full Stripping Write Mesh_Output 92.8GB



Striping – MM Level 3



MPI-IO - 7.2 B Tetra Mesh



Perspectives

BGAS (Blue Gene Active Storage) System

The Active Storage Project is aimed at:-

- -enabling close integration of emerging solid-state storage technologies with high performance networks and integrated processing capability
- -exploring the application and middleware opportunities presented by such systems
- -anticipating future scalable systems comprised of very dense Storage Class Memories (SCM) with fully integrated processing and network capability

Project to test Code_Saturne on the BGAS System (Collaboration between STFC (the Hartree Centre) and IBM)



THANK YOU FOR YOUR ATTENTION